

## THE IMPLEMENTATION OF A SPIHT WAVELET CODEC

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**Abstract.** Finding an efficient way to encode the subband coefficients is essential for revealing the potential of the wavelet transform. Shapiro's paper [1] introduces the zerotrees, as an instrument for taking advantage of the correlation between the coefficients in several subbands. The SPIHT codec presented in [2] is also based on zerotrees, but achieves better compression. In this paper are presented the results of implementing a Java codec based on the SPIHT algorithm, with little optimisations.

MSC: 68P30

Keywords: image compression

### 1. Zerotrees

A zerotree is a quad-tree in which all the nodes are smaller or equal with the root. Such a tree can be coded with a single symbol and reconstructed by the decoder as a quad-tree filled with zeros.

The coefficients obtained by applying the nonstandard wavelet transform [3] to an image can be represented with trees, because of the subsampling that is performed in the transform. Each coefficient in a low subband has four descendents in the next higher subband, and every descendent has four other descendents in the next higher subband.

Natural images usually have a low pass spectrum. The energy in the subbands decreases as the scale decreases, so the wavelet coefficients will be smaller in the higher subbands than in the lower subbands [4], [5]. Fig. 1 presents the first  $16 \times 16$  wavelet coefficients, (the first four subbands) for the Lena portrait, generated with the Villasenor 18/15 filter [6].

### 2. The SPIHT Codec

The SPIHT encoder exploits the zerotrees, based on the property that the wavelet coefficients usually decrease with scale, and there is a good probability that all the coefficients in a quad-tree will be smaller than a certain threshold, if the root is smaller than this threshold. If the image is scanned going from the lowest subband to higher subbands, many positions can be coded efficiently with zerotrees symbols.

The algorithm builds the following data structures: The *List of Insignificant Sets* LIS that contains the roots of the zerotrees, and two lists that hold the coefficients extracted from the transformed image: the *List of Insignificant Pixels* LIP, and the *List of Significant Pixels* LSP. The LIP keeps the coefficients, that are lower than a certain threshold, and the LSP holds the coefficients that are bigger than the threshold [2]. The following sets of coordinates are used for presenting the SPIHT algorithm:

- $D(i,j)$  = all the descendents of the node  $(i,j)$  in the quad-tree
- $O(i,j) = ((2i,2j), (2i,2j+1), (2i+1,2j), (2i+1,2j+1))$
- $L(i,j) = D(i,j) - O(i,j)$

$$S_n(i,j) = \begin{cases} 1, & \text{if } \max_{i,j} |c_{i,j}| \geq n \\ 0, & \text{otherwise} \end{cases}$$

The LIS contains elements  $(i,j)$  of type A that represent  $D(i,j)$ , and elements of type B, that represent  $L(i,j)$ .

### 2.1. The encoding algorithm

The image has  $l \times l$  pixels, and  $l$  is a power of 2.  $c_{i,j}$ ,  $0 \leq i < l$ ,  $0 \leq j < l$ , are the wavelet transform coefficients.

The LIP contains the coefficients  $c_{i,j}$  that are below the threshold  $n$ , and LSP the coefficients that exceeded the threshold, in absolute value. The elements of the LIS have the following structure:  $(i,j,type)$ .

The SPIHT encoder exploits the zerotree based on the property that the wavelet coefficients usually decrease with scale, and there is a good probability that the root of a quad-tree will be smaller than a certain threshold, if the root is smaller than the threshold, the image is scanned from the lowest subband with zerotree symbols.

- output  $n = 2^{\ell_{n+1}}$  and the overall average  $c_{i,j}$
  - initialize LSP as an empty list
  - add the coefficients  $c_{i,j}$  in the lowest subband to the LIP
  - add the coordinates  $(i,j)$  in the lowest subband to the LIS, as type A entries
2. for each coefficient  $c_i$  in the LIP do:

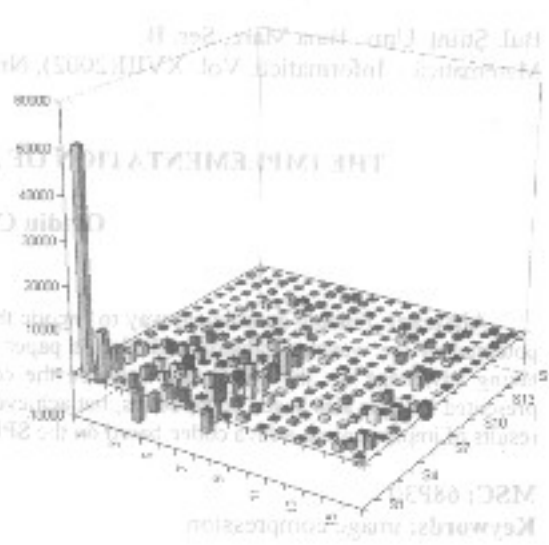


Figure 1. Coefficients of the wavelet transform.

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if  $|c_x| \geq n$ , then
    output 1
    add the value  $|c_x| - n$  to the LSP
    output the sign of the coefficient  $c_x$ 
    remove the coefficient  $c_x$  from the LIP
else output 0
for each entry  $(i,j)$  in the LIS do:
    if the entry is of type A, then
        output  $S_n(D(i,j))$ 
        if  $S_n(D(i,j)) = 1$ , then
            for each  $(k,l) \in O(i,j)$  do:
                if  $|c_{kl}| \geq n$ , then
                    output 1
                    add the value  $|c_{kl}| - n$  to the LSP
                    output the sign of the coefficient  $c_{kl}$ 
                else
                    output 0
                    add the coefficient  $c_{kl}$  to the LIP
            if  $L(i,j) \neq \phi$ , then
                set the type B for the element  $(i,j)$  in the LIS
                encode_type_B_element( $i,j$ )
            else
                remove the element  $(i,j)$  from the LIS
        else
            encode_type_B_element( $i,j$ )
3. for each element  $c_x$  in the LSP, excepting the ones included in the last pass, do:
    if  $c_x \geq n$ , then:
        output 1,  $c_x = c_x - n$ 
    else output 0
4.  $n = n/2$ , go to step 2.

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The condition for ending the process, can be related to the reach of a certain value for the current treshold  $n$ , or to the length in bits of the resulting code.

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encode_type_B_element( $i,j$ ):
output  $S_n(L(i,j))$ 
if  $S_n(L(i,j)) = 1$ , then
    add each element  $(k,l) \in O(i,j)$  to the end of LIS, as a type A element.
    remove the element  $(i,j)$  from the LIS

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## 2.2. The decoding algorithm

The elements of the LSP have the structure:  $(i, j, c_{i,j})$ , those of the LIP have the structure:  $(i, j)$ , and those of the LIS have the structure:  $(i, j, \text{type})$ .

1. read  $n, l$  and the overall average  $c_{0,0}$   
 add the coordinates of the lowest subband to the LIS, as elements of the form  $(i, j, A)$   
 add the elements  $(i, j)$  in the lowest subband to the la LIP  
 initialize the LSP, as an empty list
2. for each element  $(i, j)$  in the LIP do:
  - read the next bit of the coefficient  $c_{i,j}$ . The coefficients are read starting with the MSB.
  - if the read bit = 1, then
    - read the sign  $s_{i,j}$  of the coefficient  $c_{i,j}$
    - add the element  $(i, j, s_{i,j} \cdot n)$  to the LSP
    - remove the element  $(i, j)$  from the LIP
  - for each element  $(i, j)$  in the LIS do:
    - if the element  $(i, j)$  is of type A, then
      - read  $S_n(D(i, j))$
      - if  $S_n(D(i, j)) = 1$ , then
        - for each  $(k, l) \in O(i, j)$  do:
          - read the next bit of the coefficient  $c_{k,l}$
          - if the read bit = 1, then
            - read the sign  $s_{k,l}$  of the coefficient  $c_{k,l}$
            - add the element  $(k, l, s_{k,l} \cdot n)$  to the LSP
          - else
            - add the element  $(k, l)$  to the LIP
        - if  $L(i, j) \neq \phi$ , then
          - set the type B for the element  $(i, j)$  in the LIS
          - decode\_type\_B\_element( $i, j$ )
        - else
          - remove the element  $(i, j)$  from the LIS
      - else
        - decode\_type\_B\_element( $i, j$ )
3. for each element  $(i, j, c_{i,j})$  in the LSP, excepting the ones added in the last pass, do:
  - read the next bit of the coefficient  $c_{i,j}$
  - if the read bit = 1, then
    - $c_{i,j} = c_{i,j} + s_{i,j} \cdot n$ , where  $s_{i,j}$  is the sign of the coefficient  $c_{i,j}$
4.  $n = n / 2$ , go to step 2.

The process can continue until it is reached a certain threshold  $n$ , or a certain compression ratio, or to the end of the input data. In the end of the algorithm, the reconstructed coefficients are in the LSP, in the form of  $(i,j,c_{ij})$  elements.

decode\_type\_B\_element( $i,j$ ):

read  $S_e(L(i,j))$

if  $S_e(L(i,j)) = 1$ , then

- add each element  $(k,l)$  in  $O(i,j)$  to the LIS, as a type A element
- remove  $(i,j)$  from the LIS

### 3. Implementation results

My implementation follows the SPIHT algorithm, with a few optimisations. The next two images present the performance of the original SPIHT algorithm, compared to my implementation (MTE4).

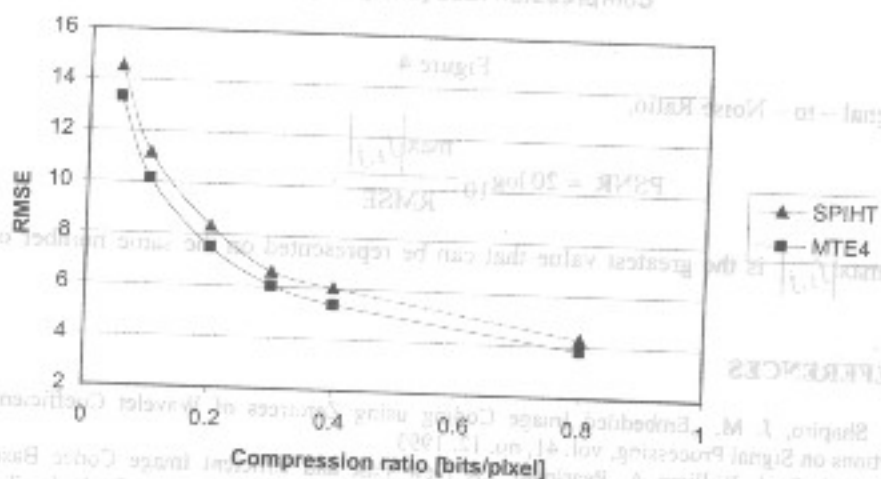


Figure 3

Root Mean Square Error:

$$RMSE = \frac{1}{M} \sqrt{\sum_{i=0}^{M-1} \sum_{j=0}^{M-1} (f_{i,j} - \tilde{f}_{i,j})^2}$$

where  $f_{i,j}$  and  $\tilde{f}_{i,j}$  are the pixels of the original, and those of the reconstructed image.

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Received: 19.09.2002

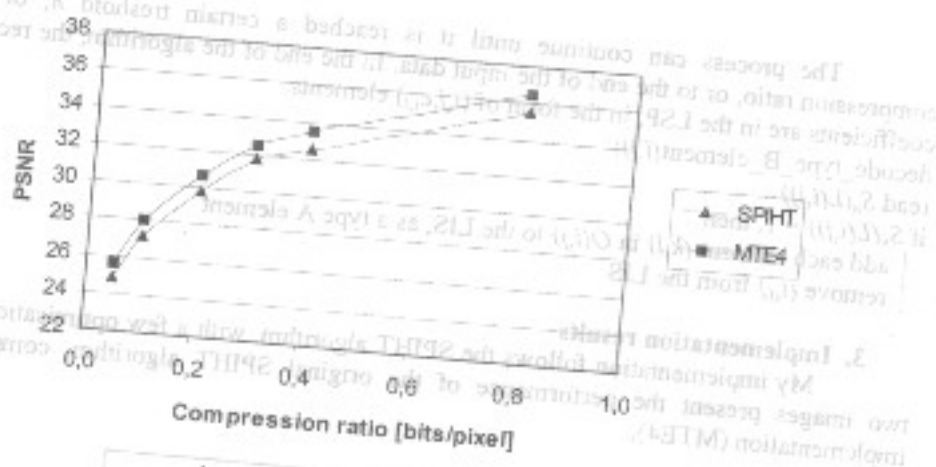


Figure 4

Peak Signal - to - Noise Ratio,

$$PSNR = 20 \log_{10} \frac{\max |f_{i,j}|}{RMSE}$$

where  $\max |f_{i,j}|$  is the greatest value that can be represented on the same number of bits as  $f_{i,j}$ .

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Received: 19.09.2002

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